

Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application.

Listing of Claims:

1. **(Currently Amended)** An apparatus for allowing a user to model at least one aspect variation of a software artifact by using extension types ~~said apparatus comprising a processor and a memory storing code accessible by the processor to provide extension types, comprising:~~

a processor;

a memory accessible by the processor;

instructions contained in the memory and executable by the processor for allowing a user to obtain a controllable software artifact, said instructions comprising:

instructions for providing a user with a software artifact; and

instructions for allowing the user to add features to the software artifact;

wherein the features comprise extension types, each extension type comprising an ordered tuple of a plurality of element types, each of the element types corresponding to different class hierarchies; and

wherein said extension types are utilized to ~~implement~~ simplify implementation of data classifications.

2. **(Original)** The apparatus according to Claim 1, wherein each extension type comprises an extension or variation of element types.

3. **(Original)** The apparatus according to Claim 1, wherein said extension types are adapted to compose classes horizontally.

4. **(Original)** The apparatus according to Claim 1, wherein each extension type is adapted to masquerade as any associated element type.

5. **(Original)** The apparatus according to Claim 1, wherein each extension type is a subtype of its associated element types.

6. **(Original)** The apparatus according to Claim 1, wherein:

each extension type has a size corresponding to the number of elements associated with the extension type; and

given two extension types α and β , a sub-type relation $\alpha <: \beta$ is definable as follows:

$|\alpha| \geq |\beta|$; and

$\alpha(0) <: \beta(0), \alpha(1) <: \beta(1), \dots, \alpha(|\beta|-1) <: \beta(|\beta|-1).$

7. **(Original)** The apparatus according to Claim 1, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta < \alpha$:

a method dispatch $p.m$ comprises starting at the element type $\beta(0)$ and walking up the class hierarchy of $\beta(0)$ to find the closest m , wherein if m is not defined in the class hierarchy of $\beta(0)$, then m is sought in the $\beta(1)$ class hierarchy and, if needed, in one or more iteratively successive class hierarchies, until found.

8. **(Original)** The apparatus according to Claim 1, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta < \alpha$:

a method dispatch p^*m comprises, for each element type $\beta(i)$, in the order $i=0, \dots, |\beta|-1$, walking up the class hierarchy of $\beta(i)$ to find the closest m in $\hat{\beta}(i)$ and dispatching the method m (if found), whereby a type error arises if m is not defined in at least one of the class hierarchies $\hat{\beta}(i)$, $i=0, \dots, |\beta|-1$.

9. **(Original)** The apparatus according to Claim 1, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta < \alpha$:

a method dispatch $p(1,3,4).m$ comprises reviewing only a class hierarchy of $\Downarrow(1)$, $\Downarrow(3)$, and $\Downarrow(4)$ to find the closest m , wherein a type error arises if m is not defined in any of $\Downarrow(1)$, $\Downarrow(3)$, or $\Downarrow(4)$.

10. **(Original)** The apparatus according to Claim 1, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta < \alpha$:

a method dispatch $p(1,3,4).m$ comprises reviewing only a class hierarchy of $\Downarrow(1)$, $\Downarrow(3)$, and $\Downarrow(4)$ to find the closest m in $\Downarrow(i)$ and dispatching the method m if found, whereby a type error arises if in any of the class hierarchies to which $\Downarrow(1)$, $\Downarrow(3)$, or $\Downarrow(4)$ belongs m is not defined.

11. **(Currently Amended)** A computer implemented method for allowing a user to model ~~of modeling~~ at least one aspect variation of a software artifact by using extension types, said method comprising:

providing the user with a controllable software artifact, said providing comprising:

providing the user with a software artifact; and

allowing the user to add features to the software artifact, wherein the features comprise ~~the step of providing~~ extension types, each extension type comprising an ordered tuple of a plurality of element types, each of the element

types corresponding to different class hierarchies, wherein said extension types are stored in a memory of at least one general-purpose computer; and

wherein said extension types are utilized to implement ~~simplify implementation of~~ data classifications.

12. **(Original)** The method according to Claim 11, wherein each extension type comprises an extension or variation of element types.

13. **(Original)** The method according to Claim 11, wherein the extension types are adapted to compose classes horizontally.

14. **(Original)** The method according to Claim 11, wherein each extension type is adapted to masquerade as any associated element type.

15. **(Original)** The method according to Claim 11, wherein each extension type is a subtype of its associated element types.

16. **(Original)** The method according to Claim 11, wherein:

each extension type has a size corresponding to the number of elements associated with the extension type; and

given two extension types α and β , a sub-type relation $\alpha <: \beta$ is definable as follows:

$$|\alpha| \geq |\beta|; \text{ and}$$

$$\alpha(0) <: \beta(0), \alpha(1) <: \beta(1), \dots, \alpha(|\beta|-1) <: \beta(|\beta|-1).$$

17. **(Original)** The method according to Claim 11, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta <: \alpha$:

a method dispatch $p.m$ comprises starting at the element type $\beta(0)$ and walking up the class hierarchy of $\beta(0)$ to find the closest m , wherein if m is not defined in the class hierarchy of $\beta(0)$, then m is sought in the $\beta(1)$ class hierarchy and, if needed, in one or more iteratively successive class hierarchies, until found.

18. **(Original)** The method according to Claim 11, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta <: \alpha$:

a method dispatch $p*m$ comprises, for each element type $\beta(i)$, in the order $i=0, \dots, |\beta|-1$, walking up the class hierarchy of $\beta(i)$ to find the closest m in $\hat{\Downarrow}(i)$ and dispatching the method m (if found), whereby a type error arises if m is not defined in at least one of the class hierarchies $\hat{\Downarrow}(i)$, $i=0, \dots, |\beta|-1$.

19. **(Original)** The method according to Claim 11, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta <: \alpha$:

a method dispatch $p(1,3,4).m$ comprises reviewing only a class hierarchy of $\Downarrow(1)$, $\Downarrow(3)$, and $\Downarrow(4)$ to find the closest m , wherein a type error arises if m is not defined in any of $\Downarrow(1)$, $\Downarrow(3)$, or $\Downarrow(4)$.

20. **(Original)** The method according to Claim 11, wherein, with α being the extension type of a variable p and β being the runtime extension type of the object pointed by p , so that $\beta <: \alpha$:

a method dispatch $p(1,3,4).m$ comprises reviewing only a class hierarchy of $\Downarrow(1)$, $\Downarrow(3)$, and $\Downarrow(4)$ to find the closest m in $\Downarrow(i)$ and dispatching the method m if found, whereby a type error arises if in any of the class hierarchies to which $\Downarrow(1)$, $\Downarrow(3)$, or $\Downarrow(4)$ belongs m is not defined.

21. **(Currently Amended)** A program storage device readable by machine, tangibly encoded with a program of instructions executable by a processor of the machine to perform method steps for allowing a user to model at least one variation of a software artifact by using extension types, said method steps comprising:

A data storage device readable by machine, comprising a data structure stored on the device, the data structure

providing the user with a controllable software artifact, said providing comprising:

providing the user with a software artifact; and

allowing the user to add features to the software artifact, wherein the features comprise ~~being~~ at least one extension type comprising an ordered tuple of a plurality of element types, each of the element types corresponding to different class hierarchies;

wherein said at least one extension type allows a user to ~~model at least one aspect of a software artifact to simplify implementation of~~ implement data classifications.